

Efficient Reduced-Order Modeling of Frequency-Dependent Coupling Inductances Associated with 3-D Interconnect Structures

Luís Miguel Silveira, *Member, IEEE*, Mattan Kamon, *Student Member, IEEE*, and Jacob White, *Member, IEEE*

Abstract—Reduced-order modeling techniques are now commonly used to efficiently simulate circuits combined with interconnect, but generating reduced-order models from realistic three-dimensional (3-D) structures has received less attention. In this paper, we describe a Krylov-subspace based method for deriving reduced-order models directly from the 3-D magneto-quasistatic analysis program FastHenry. This new approach is no more expensive than computing an impedance matrix at a single frequency.

Index Terms—Interconnect, Krylov-subspace, inductance, reduced order model, packaging analysis, computational electromagnetics.

I. INTRODUCTION

THE dense three-dimensional packaging used in compact electronic systems may produce magnetic interactions which interfere with system performance. Such effects are difficult to simulate because they occur only as a result of an interaction between the field distribution in a complicated geometry of conductors, and the circuitry connected to those conductors. Recent work on reduced-order modeling techniques has made it possible to efficiently simulate circuits combined with interconnect [1], but generating the reduced-order models from realistic three-dimensional (3-D) structures has received less attention. The most commonly used approach to generating reduced-order models is to use a 3-D field solver to compute impedance matrices over a range of frequencies, and then use a rational function fitting algorithm [2]. This approach has been shown to produce accurate frequency-domain reduced-order models which are easily included in a standard circuit simulator [3].

In order to use frequency-domain fitting as described above, it is necessary to use the field solver to compute impedance

matrices at dozens of frequency points, and this is computationally expensive. It is possible to derive a more efficient approach by exploiting the fact that 3-D field solvers typically use Krylov-subspace based iterative methods. These iterative methods can provide more than just a solution at a particular frequency, they can be used to directly construct reduced-order models [4].

In this paper, we present a numerically robust and accurate approach for computing reduced-order models of magneto-quasistatic coupling in complicated 3-D structures. The approach is based on using the multipole-accelerated program FastHenry [5], combined with the Krylov-subspace based Arnoldi algorithm [6]. We begin, in Section II, by describing the mesh-formulation approach of FastHenry. In Section III, the standard Padé approximation approach as well as an Arnoldi-based approach are derived. In Section IV, results are presented comparing the accuracy of the two model-order reduction methods on an RLC filter and a package example. Finally, in Section V, we present conclusions and acknowledgments.

II. THE MESH FORMULATION APPROACH

The frequency dependent resistance and inductance matrices describing the terminal behavior of a set of conductors can be rapidly computed with the multipole-accelerated mesh-formulation approach as implemented in FastHenry [5]. To describe the approach, consider that each conductor is approximated as piecewise-straight sections. The volume of each straight section is then discretized into a collection of parallel thin filaments through which current is assumed to flow uniformly.

To derive a system of equations for the filament currents, we start by assuming the system is in sinusoidal steady-state, and following the partial inductance approach in [7], the branch current phasors can be related to branch voltage phasors by

$$\begin{aligned} V_b &= (R + j\omega L)I_b \\ &= ZI_b \end{aligned} \quad (1)$$

where $V_b, I_b \in \mathbb{C}^b$, b is the number of branches (number of current filaments), and ω is the excitation frequency. The entries of the diagonal matrix $R \in \mathbb{R}^{b \times b}$ represent the DC resistance of each current filament, and $L \in \mathbb{R}^{b \times b}$ is the dense matrix of partial inductances.

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L. M. Silveira is with the Electronics Division of the Instituto de Engenharia de Sistemas e Computadores (INESC), Lisbon, Portugal.

M. Kamon and J. White are with the Department of Electrical Engineering and Computer Science, Research Laboratory of Electronics, Massachusetts Institute of Technology, Cambridge, MA 02139 USA.

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Kirchhoff's voltage law, which implies that the sum of branch voltages around each mesh (a mesh is any loop of branches in the graph which does not enclose any other branches) in the network is represented by

$$\begin{aligned} MV_b &= V_s \\ M^T I_m &= I_b \end{aligned} \quad (2)$$

where $V_s \in \mathbb{C}^m$ is the mostly zero vector of source branch voltages, $I_m \in \mathbb{C}^m$ is the vector of mesh currents, $M \in \mathbb{R}^{m \times b}$ is the mesh matrix, and m is the number of meshes, which is typically somewhat less than b , the number of filaments. Combining (1) and (2) yields

$$MZM^T I_m = V_s. \quad (3)$$

The complex admittance matrix which describes the external terminal behavior of a t -conductor system, denoted $Y_t = Z_t^{-1}$, can be derived from (3) by noting that

$$I_t = Y_t V_t.$$

I_t and V_t are the terminal source currents and voltages of the t -conductor system, which are related to the mesh quantities by $I_t = N^T I_m$, $V_s = N V_t$, where $N \in \mathbb{R}^{m \times t}$ is a terminal incidence matrix determined by the mesh formulation.

Hence, to compute the i th column of Y_t , solve (3) with a V_s whose only nonzero entry corresponds to V_{t_i} , and then extract the entries of I_m associated with the source branches.

To solve (3) by Gaussian elimination would require $\mathcal{O}(m^3)$ operations where the number of meshes, m , can exceed 10 000 when modeling frequency dependent effects. Instead, programs like FastHenry solve (3) using a multipole-accelerated GMRES iterative algorithm [6], which requires $\mathcal{O}(b)$ operations. The complexity is reduced from $\mathcal{O}(m^3)$ to $\mathcal{O}(m^2)$ by using GMRES instead of Gaussian elimination, and then to $\mathcal{O}(b)$ by using a hierarchical multipole algorithm [8].

III. REDUCED-ORDER MODELING

One approach to coupling package models with circuits is to simply include a sparse tableau version of (3) in a circuit simulator [9]. However, since MZM^T is large and dense, this approach can be computationally intractable. A more computationally efficient approach is to represent (3) with a reduced-order model.

A. State-Space Formulation

As mentioned in the introduction, to use frequency-domain fitting to generate a reduced-order model for the frequency-dependent entries of Y_t , it would be necessary to construct and solve (3) for dozens of values of ω . To derive a more efficient approach, consider forming the state-space representation of (3). To that end, expand Z into $R + sL$ to get

$$\begin{aligned} s(MLM^T)I_m &= -(MRM^T)I_m + N V_t \\ I_t &= N^T I_m. \end{aligned} \quad (4)$$

With the representation in (4), the (i, j) th entry of the complex admittance matrix can be computed using a set of terminal voltages whose only nonzero entry corresponds to V_{t_j} , and can be written as

$$\begin{aligned} \frac{I_{t_i}}{V_{t_j}} &= Y_{t_{ij}}(s) \\ &= c^T (I - sA)^{-1} b \end{aligned} \quad (5)$$

where $A = -(MRM^T)^{-1}(MLM^T)$ and $b = (MRM^T)^{-1}N_j$ and $c = N_i$, where N_i indicates the i th column of N . It is possible to derive extensions of all the methods mentioned in this paper to directly compute approximations to the full t input, t output system in (4) [10], [11]. In the remainder, however, we will restrict our discussion to single-input single-output systems characterized by a transfer function such as (5).

The standard approach to computing a reduced-order model of (5) is to use Padé approximations [12]. Padé approximates are rational functions in the Laplace transform variable s which match power series expansions of the original system transfer function up to some order. More specifically, consider the admittance

$$\begin{aligned} Y_{t_{ij}}(s) &= c^T (I - sA)^{-1} b \\ &= \sum_{k=0}^{\infty} m_k s^k \end{aligned} \quad (6)$$

where

$$m_k = c^T A^k b \quad (7)$$

is the k th moment of the transfer function. A Padé approximation of q th order is defined as the rational function

$$G_q^P(s) = \frac{b_{q-1}s^{q-1} + \dots + b_1s + b_0}{a_q s^q + a_{q-1}s^{q-1} + \dots + a_1s + 1} \quad (8)$$

whose coefficients are selected to match the first $2q - 1$ moments of the transfer function (5).

Padé approximates can be computed using direct evaluation of the moments, though the approach is ill-conditioned because such computation relies on a power iteration with the system matrix A . Instead, Lanczos-style algorithms can be used that are numerically more robust [4].

B. Arnoldi-Based Approximations

An alternative approach, which robustly generates a somewhat different approximation than Padé, can be derived using an Arnoldi process as in the GMRES algorithm. The idea behind this approach is similar to that of [4], and is that of selecting an orthonormal basis for the Krylov subspace $\mathcal{K}_k(A, b) = \text{span}\{b, Ab, A^2b, \dots, A^{k-1}b\}$ using a modified Gram-Schmidt process. See the basic outline of the Arnoldi algorithm given in Algorithm 1.

Algorithm 1. (Arnoldi Process):

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arnoldi(in:  $\mathbf{A}$ ,  $\mathbf{b}$ ,  $q$ ; out:  $\mathbf{V}_q$ ,  $\mathbf{v}_{q+1}$ ,  $\mathbf{H}_q$ ,  $h_{j+1,j}$ )
{
   $\mathbf{v}_1 = \mathbf{b}/\|\mathbf{b}\|_2$ 
  for ( $j = 1$ ;  $j \leq q$ ;  $j++$ ) {
     $\mathbf{w} = \mathbf{A}\mathbf{v}_j$ 
    for ( $i = 1$ ;  $i \leq j - 1$ ;  $i++$ ) {
       $h_{i,j} = \mathbf{w}^T \mathbf{v}_i$ 
       $\mathbf{w} = \mathbf{w} - h_{i,j} \mathbf{v}_i$ 
    }
     $h_{j+1,j} = \|\mathbf{w}\|_2$ 
    if ( $h_{j+1,j} \neq 0$ ) {
       $\mathbf{v}_{j+1} = \mathbf{w}/h_{j+1,j}$ 
    }
  }
   $\mathbf{V}_q = [\mathbf{v}_1 \cdots \mathbf{v}_q]$ 
   $\mathbf{H}_q = (h_{i,j}), \quad i, j = 1, \dots, q$ 
}

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Three aspects of Algorithm 1 should be noted. First, the orthogonality between the basis vectors makes the Arnoldi algorithm a better conditioned process than direct evaluation of the moments. Second, the computation of \mathbf{b} is inexpensive since \mathbf{MRM}^T is sparse. Finally, because L is dense, the dominant cost of each step of an Arnoldi process is a matrix-vector product, $\mathbf{A}\mathbf{x} = -(\mathbf{MRM}^T)^{-1}(\mathbf{MLM}^T)\mathbf{x}$. In practice, the matrix-vector product cost dominates over the inner loop orthogonalization even when the dense part, $(\mathbf{MLM}^T)\mathbf{x}$, is rapidly computed with a hierarchical multipole-algorithm as in FastHenry.

After q steps, the Arnoldi algorithm returns a set of q orthonormal vectors, as the columns of the matrix $\mathbf{V}_q \in \mathbb{R}^{m \times q}$, and a $q \times q$ upper Hessenberg matrix \mathbf{H}_q whose entries are the Gram-Schmidt orthogonalization coefficients $h_{i,j}$. Following the approach in [4], these two matrices satisfy the following relationship:

$$\mathbf{A}\mathbf{V}_q = \mathbf{V}_q\mathbf{H}_q + h_{j+1,j}\mathbf{v}_{q+1}\mathbf{e}_q^T \quad (9)$$

where \mathbf{e}_q is the q th unit vector in $\mathbb{R}^{m \times m}$.

From (9), it can easily be seen that after q steps of an Arnoldi process, for $k < q - 1$

$$\begin{aligned} \mathbf{A}^k \mathbf{b} &= \|\mathbf{b}\|_2 \mathbf{A}^k \mathbf{V}_q \mathbf{e}_1 \\ &= \|\mathbf{b}\|_2 \mathbf{V}_q \mathbf{H}_q^k \mathbf{e}_1. \end{aligned} \quad (10)$$

With this relation, the moments can be related to \mathbf{H}_q by

$$\begin{aligned} m_k &= \mathbf{c}^T \mathbf{A}^k \mathbf{b} \\ &= \|\mathbf{b}\|_2 \mathbf{c}^T \mathbf{V}_q \mathbf{H}_q^k \mathbf{e}_1 \end{aligned} \quad (11)$$

and so the q th-order Arnoldi-based approximation to \mathbf{Y}_{ij} can be written as

$$\mathbf{G}_q^A(s) = \|\mathbf{b}\|_2 \mathbf{c}^T \mathbf{V}_q (\mathbf{I} - s\mathbf{H}_q)^{-1} \mathbf{e}_1. \quad (12)$$

Using the eigendecomposition $\mathbf{H}_q = \mathbf{S}_q \Lambda_q \mathbf{S}_q^{-1}$, the expression for the approximating rational function becomes

$$\mathbf{G}_q^A(s) = \|\mathbf{b}\|_2 \sum_{k=0}^q \frac{\mu_k \nu_k}{s - p_k} \quad (13)$$

where $\boldsymbol{\mu} = -\mathbf{c}^T \mathbf{V}_q \mathbf{S}_q \Lambda_q^{-1}$, $\boldsymbol{\nu} = \mathbf{S}_q^{-1} \mathbf{e}_1$, and $\mathbf{p} = \text{diag}(\Lambda_q^{-1})$ are the poles of the approximation.

Note that the rational function $\mathbf{G}_q^A(s)$ is *not* a Padé approximation as it has q poles, but only matches $q - 2$ moments since (10) is only valid for $k < q - 1$. However, computing the rational function requires only q matrix-vector products, roughly half the number of matrix-vector products required to compute a q th-order Padé approximate which matches $2q - 1$ moments. For the same computational effort required to compute the q th-order Padé approximant $\mathbf{G}_q^P(s)$ one could obtain $\mathbf{G}_{2q}^A(s)$, which has $2q$ poles and matches $2q - 2$ moments. We have observed empirically that the extra number of poles improves the accuracy of the Arnoldi approximation, however, the larger number of poles makes it more expensive to include in a circuit simulator.

C. Other Work

It is worth noting that a block Arnoldi algorithm can be used to more efficiently handle multiple input and output problems. Such an approach is described in [11].

Also, it can be shown that given a symmetric, stable system, \mathbf{A} , the Arnoldi approach is guaranteed to generate a stable reduced-order model, \mathbf{H}_q [13], [14]. Additionally, the inductance system, \mathbf{A} , of (4), is stable and can be made symmetric and stable with a coordinate transformation thus guaranteeing the stability of the models presented in this paper [14].

IV. EXPERIMENTAL RESULTS

In the preceding section we described algorithms to compute Padé approximations of order q and Arnoldi-based models of orders q and $2q$. In this section we compare the accuracy of these three approximations: first, for a difficult to model RLC filter example, and then when used to obtain reduced-order models for the frequency-dependent admittance for a small set of package pins. This reduced-order model is then used to investigate crosstalk between the package pins.

A. Filter Example

Figure 1 shows the Bode plots of the seventh-order Padé and the seventh and 14th-order Arnoldi-based approximations to a 14th-order RLC filter's transfer function. Also shown in the picture is the exact transfer function. For the low frequency range, all approximations are indistinguishable. However, for higher frequencies, as is clear from the figure, the seventh-order Padé and the seventh-order Arnoldi-based approximation have comparable accuracy, while the 14th-order Arnoldi-based approximation, which requires the same number of matrix-vector products as the seventh-order Padé, is indistinguishable from the exact transfer function. This last observation is not surprising, due to the finite termination properties of the algorithm. It should be noted, however, that any 14th-order approximation will be significantly more expensive to use in a circuit simulator than a seventh-order approximation. Nevertheless, the ability to compute higher orders of approximation at no extra cost remains a valuable property of the Arnoldi-based approximation method.

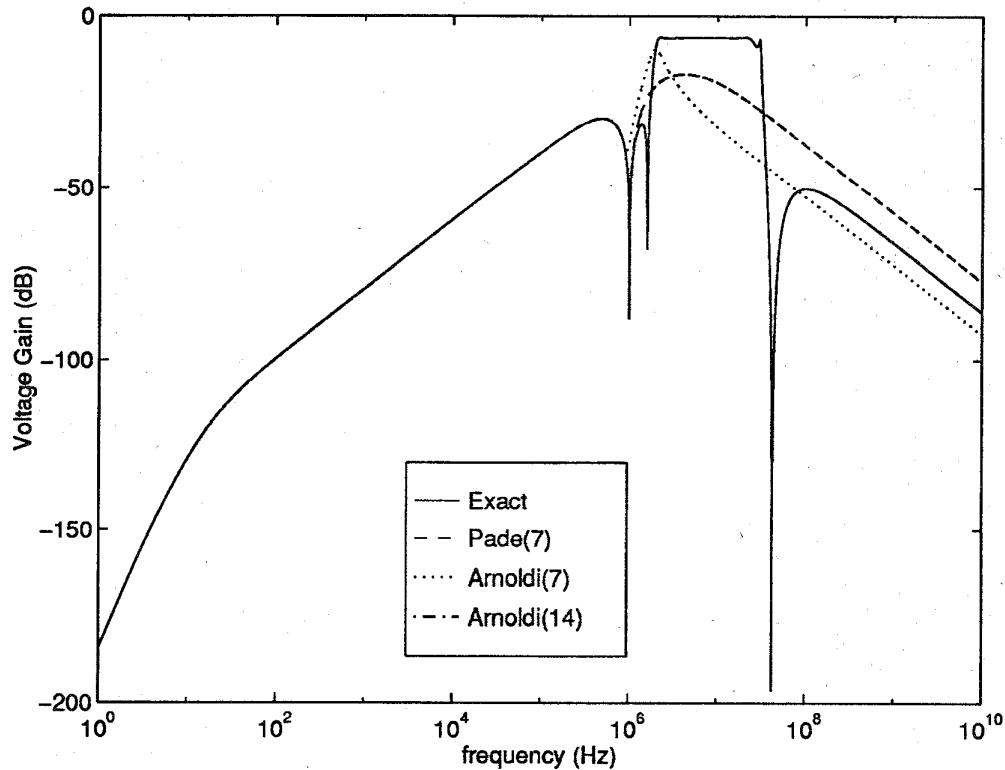


Fig. 1. Bode plots for the approximations $G_7^P(s)$, $G_7^A(s)$, and $G_{14}^A(s)$ to the RLC filter's transfer function. The exact transfer function and the 14th-order approximation are indistinguishable.

B. Package Example

Consider the small set of package pins shown in Fig. 2. To compute the resistance and inductance matrices with FASTHENRY, the pins were discretized into three filaments along their height and four along their width producing a system of size $m = 887$. This allows modeling of changes in resistance and inductance due to skin and proximity effects. Fig. 3 shows the magnitude of the error of the eighth-order Padé and the eighth and 16th-order Arnoldi-based approximations to $Y_{t_{12}}(s)$, the coupled admittance transfer function for the PEEC discretization between pins 1 and 2 computed by (5). As can be seen from the plot, all three approximation have an error well below 5%. It is worth noting that computing the eighth-order Arnoldi-based model for one column of the admittance matrix requires eight matrix-vector products, while for a rational function fitting algorithm, computing one column of the admittance matrix over a range of 50 discrete frequencies would require 140 such products.

To investigate the crosstalk effects between the package pins in Fig. 2, the configuration shown in Fig. 4 is used where it was assumed that the five middle lines carry output signals from the chip and the two outer pins carry power and ground. The signals are driven and received with CMOS inverters which are capable of driving a large current to compensate for the impedance of the package pins. The capacitance of the pin itself is assumed to be 8 pF and the interconnect from the end of pin to the receiver is modeled with a capacitance of 5

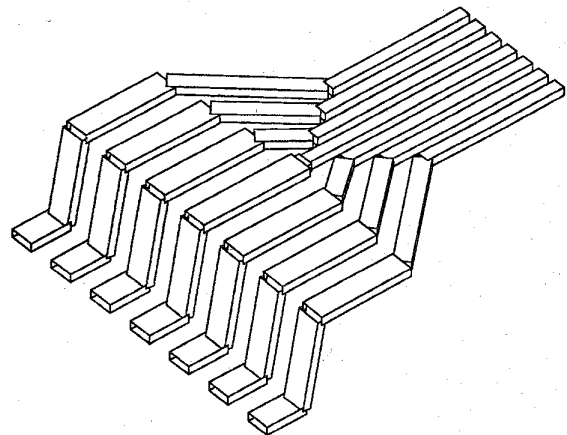


Fig. 2. Seven pins of a cerquad pin package.

pF. A 0.1 μF decoupling capacitor is connected between the driver's power and ground to minimize supply fluctuations. The frequency dependence of each element in the admittance matrix is modeled via Arnoldi-based approximations of eighth order. These models are then incorporated into SPICE3 as a frequency-dependent voltage-controlled current source. As a sample time domain simulation, imagine that at time $t_0 = 4$ ns the signal on pin 4 of Fig. 4 switches from high to low, and pins 2, 3, 5, and 6 switch from low to high but that due to delay on chip, pins 2, 3, 5, and 6 switch at $t_1 = 5$ ns. In

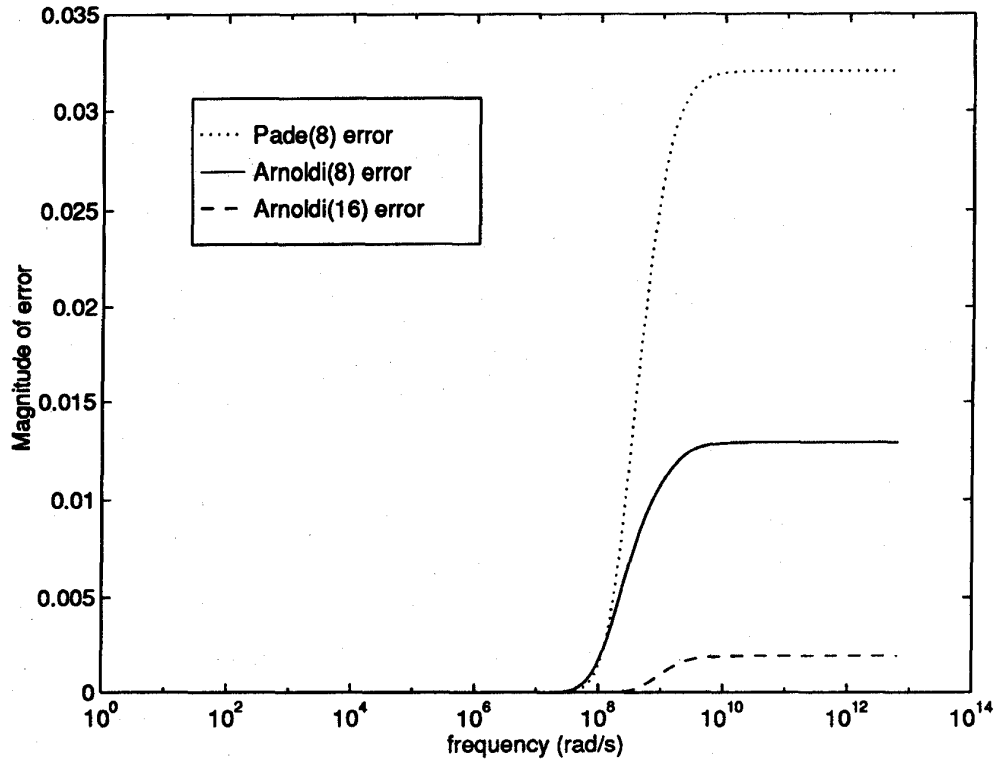


Fig. 3. Magnitude of the error for the approximations $G_8^P(s)$, $G_8^A(s)$, and $G_{16}^A(s)$ to the coupled admittance transfer function between pins 1 and 2.

this example, significant current will suddenly pass through the late pins while pin 4 is in transition. Due to crosstalk, this large transient of current has significant effects on the input of the receiver on pin 4, as shown in Fig. 5. Note that the input does not rise monotonically. Fig. 5 also shows that the bump in the waveform is carried through to the output of receiver, as a large glitch.

V. CONCLUSION AND FUTURE WORK

In this paper, we describe an accurate approach to using the iterative method in FAST HENRY to compute reduced-order models of frequency-dependent inductance matrices associated with complicated 3-D structures. The key advantage of this method is that it is no more expensive than computing the inductance matrix at a single frequency. We also compared two approaches to the model-order reduction, the reformulated Padé-based approach using the Lanczos algorithm (PVL) and an Arnoldi-based approach using an algorithm based on the Arnoldi process. We showed that the Arnoldi-based algorithm can have advantages over PVL in certain applications. In particular, in the Arnoldi-based algorithm, each set of iterations produces an entire column of the inductance matrix rather than a single entry, and if matrix-vector product costs dominate then the Arnoldi-based algorithm produces a better approximation for a given amount of computational effort.

More work is underway to understand the accuracy and robustness properties of the Arnoldi algorithm. Since the exact solution is only available in simple cases, development of error

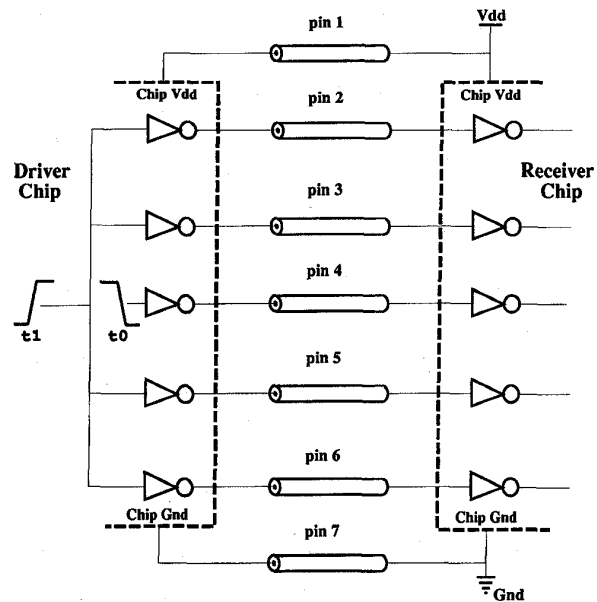


Fig. 4. General configuration for the connection between receiver and driver chips. All the circuit elements inside the same chip share that chip's power and ground.

estimates are necessary. Additionally, the algorithm can be shown to generate a guaranteed stable reduced-order model for the systems described in this paper and also any symmetric, stable system.

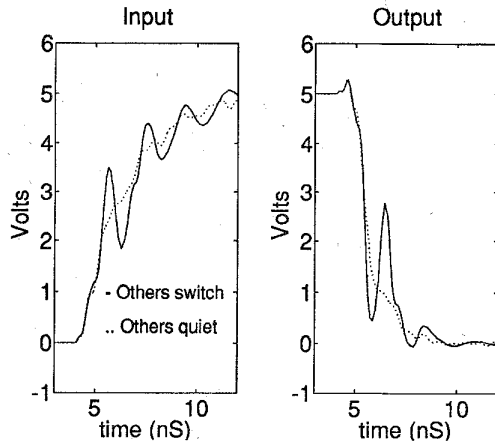


Fig. 5. Results of the timing simulation for the output of the receiver gate connected to pin 4 when the adjacent pins switch 1 ns after pin 4.

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Luís Miguel Silveira (S'85-M'95) was born in Lisbon, Portugal. He received the Engineer's (summa cum laude) and Master's degrees in electrical and computer engineering from Instituto Superior Técnico at the Technical University of Lisbon, in 1986 and 1989, respectively, and the M.S., E.E., and Ph.D. degrees from the Massachusetts Institute of Technology, Cambridge, MA, in 1990, 1991, and 1994, respectively.

His Doctoral Thesis was entitled "Model order reduction techniques for circuit simulation." From 1992 to 1994, he was supported by an IBM doctoral fellowship. He is currently an Associate Professor at *Instituto Superior Técnico of the Technical University of Lisbon*, a Researcher in the Electronics division of *INESC, Instituto de Engenharia de Sistemas e Computadores*, in Lisbon, Portugal, and a Member of the Cadence European Laboratories. His research interests are in various aspects of computer-aided design of integrated circuits with emphasis on parallel computer algorithms and the theoretical and practical issues concerning numerical simulation methods for circuit design problems.

Dr. Silveira is a member of Sigma Xi.



Mattan Kamon (S'91) received the B.S. degree in engineering science and the M.A. degree in mathematics from the Pennsylvania State University, University Park, PA, in 1991, and the S.M. degree in electrical engineering and computer science from the Massachusetts Institute of Technology, Cambridge, MA, in 1994. He is presently working toward the Ph.D. degree at the Massachusetts Institute of Technology, Cambridge, MA, with interest in developing numerical algorithms for circuit and interconnect simulation and design.



Jacob White (S'80-M'83) received the B.S. degree in electrical engineering and computer science from the Massachusetts Institute of Technology, Cambridge, MA, and the M.S. and Ph.D. degrees in electrical engineering and computer science from the University of California, Berkeley.

He worked at the IBM T. J. Watson research center from 1985 to 1987, was the Analog Devices Career Development Assistant Professor at the Massachusetts Institute of Technology from 1987 to 1989, and was a 1988 Presidential Young Investigator. He is currently an Associate Professor at MIT, and his research interests are in serial and parallel numerical algorithms for problems in circuit, interconnect, device, and microelectromechanical system design.